

Lecture 2.4]

How to DRAW BS BIMODS

①

Before: How to draw/encode morphisms. Now: how to draw/encode elements.

Choose RTG, find basis of $BS(\underline{w})$ as right R -mod. Size is $2^{\ell(w)}$. $\ell(w)$ description.

Description 1: In zero-one sequences First, silly observation: any elt of $BS(\underline{w})$ is $\Psi(10\ldots 01)$ for some $\Psi \in \text{End}_{R-\text{Bim}}(BS(\underline{w}))$ i.e. $\Psi = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$.

Description 1: OI sequences Consider Bs. We know $\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$ are right R -basis.

Better basis $G_1 = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$ canon

$$G_S = \frac{1}{2}(G_1 \otimes 1 + 1 \otimes G_1) \text{ also canon } f_{GS} = sf \begin{pmatrix} b \\ p \end{pmatrix} (101) = C_S.$$

Remember $f = \begin{pmatrix} b \\ p \end{pmatrix} sf + \begin{pmatrix} s \\ f \end{pmatrix}$ so $fog = C_S 2sf)g + G_1 sf)g$.

Consider $BS_B \dots B_n$. $\Sigma \in \{0,1\}^{\ell(w)}$

$$\Psi_\Sigma = \begin{pmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 \end{pmatrix}$$

$$\Psi_\Sigma(10\ldots 01) = C_\Sigma = G_1 C_1 \dots G_n$$

$$10\ldots 01 = C_{bot} = G_{00\dots 00}$$

$$C_{top} = G_{11\dots 11} \quad f_{C_{top}} = G_{top}f \quad f_{\begin{pmatrix} b \\ p \end{pmatrix}^d} = \begin{pmatrix} b \\ p \end{pmatrix}^d f$$

Exercise: Write f_{C_Σ} in right R basis for f linear.

$$\deg(C_\Sigma) = -d + 2|\Sigma|$$

Using $\begin{pmatrix} b \\ p \end{pmatrix} = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} + \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$, its of the form $\begin{pmatrix} b \\ p \end{pmatrix} + \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$

More properties of BSBm $BS(\underline{w})$ is a commutative ring!

$R \otimes R \otimes \dots \otimes R$ termwise mult.

grading is wrong though b/c $C_{bot} = id$ & $\deg = d$

$$\deg(a \cdot b) = \deg a + \deg b + d.$$

In OI sequences, mult = stacking, i.e. $\Psi_\Sigma(C_{bot}) \circ \Psi_\Sigma(C_{bot}) = (\Psi_\Sigma \circ \Psi_\Sigma)(C_{bot})$ (Ψ_Σ all commute)

$$\text{Ex: } \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} = b \begin{pmatrix} 0 \\ 1 \end{pmatrix} \quad \text{Ex: } \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} = 2 \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} - b \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

by poly argument

$BS(\underline{w})$ also has a trace $\text{Tr}: BS(\underline{w}) \rightarrow R$ take coeff of C_Σ (this is canonical...)

and a pairing $BS(\underline{w}) \times BS(\underline{w}) \rightarrow R$
 $\langle a, b \rangle = \text{Tr}(ab)$

$$\text{Ex: } \langle \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}, \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \rangle = 2 \quad \text{Ex: } \langle C_{top}, C_{top} \rangle = \prod \alpha_i$$

Claim: $\langle C_\Sigma, C_{1-\Sigma} \rangle = 1$ and $\langle C_\Sigma, C_{1-\Sigma} \rangle = 0$ for $\Sigma \neq 1-\Sigma$ in same order

⇒ non degenerate to DGRB6 O.

This pairing is invariant: $\deg \langle a, b \rangle = \deg a + \deg b$; $\langle af, b \rangle = \langle a, bf \rangle = \langle a, b \rangle f$ • $\langle fa, b \rangle = \langle a, fb \rangle = \langle a, f \rangle b$ (nothing)

Some abstract bullshit: $\text{Hom}_{\text{Right } R\text{-Mod}}(B, R) = \text{DB}$, it's an R -bimodule. $D(B(\underline{\omega})) = \text{DB}(\underline{\omega})$
 $\text{Hom}_{R\text{-Bim}}(B, \text{DB})^\circ = \text{space of invariant forms on } B$
 isom \longleftrightarrow nondegen to degree 0

Thus $BS(\underline{\omega})$ has nondeg form $\Rightarrow BS(\underline{\omega})$ is self-dual. What about other SBim?

Easy Exercise: R_w std bim is self-dual.

Thus SCT $\Rightarrow B_w \in BS(\underline{\omega})$ so $D B_w \oplus DBS(\underline{\omega})$, ! property $\Rightarrow B_w \cong DB_w$
 rex \hookrightarrow All indecomp SBim are self-dual $\Rightarrow \exists$ nondegen int form.

What is this form? Clearly B_w inherits a form from $BS(\underline{\omega})$.

Exercise: $C_{\text{bot}} \in B_w$ (use support filtration). Thus restricted form is nonzero.
 $C_{\text{top}} \in B_w$ (use duality)

Cor: Suppose S. Conf, so that $\text{End}(B_w) = R$. Then any ~~nonzero~~ int form is
 the unique nondegen int form (up to R^\times).

We'll prove this en route to proving S. Conf later. Example to come

Second description of elts of $BS(\underline{\omega})$, also parametrized by 2^{nd} term...

Any elt is $\Psi(C_{\text{bot}})$ for some $\Psi \in \text{End}(BS(\underline{\omega}))$. Well, $\Psi = \sum \frac{\omega}{\tau} f$.
 What is $\Psi(C_{\text{bot}}) = ?$

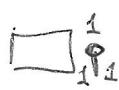
Prop: For each x appearing as a subexp. of $\underline{\omega}$, $\exists!$ subexp. can_x s.t.

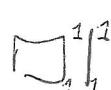
$\text{LL}_{\text{can}_x}: BS(\underline{\omega}) \rightarrow BS(x)$ sends ~~C_{bot}~~ to something nonzero!

It has maximal degree and is minimal in the Bruhat path dominance order.

Moreover, $\text{LL}_{\text{can}_x}(C_{\text{bot}, \underline{\omega}}) = C_{\text{bot}, x}$

Pf: rex more sends $C_{\text{bot}} \rightarrow C_{\text{bot}}$ so must have only ups!

If u_0  If DO  Exercise \Rightarrow the rest of the prop

If u_1  If D_1 

Rmk: LL_{can_x} depends on a choice of rex move , but two choices differ by a term which kills C_{bot} . (3)

Rmk:

$$\begin{array}{c} \text{Can} \\ \text{Can} \end{array} = \begin{array}{c} \text{Can} \end{array}.$$

So $\Psi(C_{\text{bot}}) = \sum \begin{array}{c} \text{rr} \\ \text{Can} \end{array} f(C_{\text{bot}}) = \sum \begin{array}{c} \text{rr} \\ \text{Can} \end{array}^w f(C_{\text{bot}}, x)$ finally, given by z^d terms.

upwards light leaves are injective

In this description what is multiplication $\begin{array}{c} \text{rr} \\ \text{Can} \end{array}^w \begin{array}{c} \text{rr} \\ \text{Can} \end{array}^w = ?$ Not obvious!

$$\left\langle \begin{array}{c} \text{rr} \\ \text{Can} \end{array}^w, \begin{array}{c} \text{rr} \\ \text{Can} \end{array}^w \right\rangle$$

COB to description 1?

One nice fact:

Claim: If $\Psi(C_{\text{bot}}) = 0$ then coeff of $\begin{array}{c} \text{Can} \\ \text{leaf} \end{array}$ is zero, i.e. its in lower terms \mathcal{D}^{new} .

$$\Psi = \sum \begin{array}{c} \text{rr} \\ w \end{array} \xrightarrow{\text{anyting}} \xleftarrow{\text{red exp}}$$

Pf: $\Psi(C_{\text{bot}}) = \sum (\text{coeff of } \begin{array}{c} \text{LL} \\ \text{Can} \end{array}^w) \cdot \text{symbol } \begin{array}{c} \text{LL} \\ w \end{array} \xleftarrow{\text{in indep.}}$

In particular, this coeff is zero

Will use later.

TIME FOR QUESTIONS